Axiom schema of Markov's principle preserves disjunction and existence properties

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Introduction: Disjunction and Existence Properties

"Hallmarks" of constructivity of intuitionistic logic \mathbf{H}_* :

Fact

 \mathbf{H}_* has the disjunction property (DP); for every $A \lor B$: $\mathbf{H}_* \vdash A \lor B \Rightarrow \mathbf{H}_* \vdash A$ or $\mathbf{H}_* \vdash B$.

 \mathbf{H}_* has the existence property (EP); for every $\exists xA(x): \mathbf{H}_* \vdash \exists xA(x) \Rightarrow$ there exists a v such that $\mathbf{H}_* \vdash A(v)$.

N.B. A(v) should be taken as a formula congruent to A free from collision of variables.

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H_{*} has the existence property (EP); for every $\exists xA(x)$: **H**_{*} $\vdash \exists xA(x) \Rightarrow$ there exists a *v* such that **H**_{*} $\vdash A(v)$.

 $\mathbf{H}_* + A$: the logic obtained from \mathbf{H}_* by adding the axiom schema A. There are schmemata A such that $\mathbf{H}_* + A$ enjoys both of DP and EP.

We are interested in such schemata (i.e., $\mathbf{H}_* + A$ still enjoys DP and EP) in the setting of Intermediate Predicate Logics, particularly in those schemata related to constructive theories.

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Markov's Principle and Limited Principle of Omniscience

In the setting of intermediate Predicate Logics, we consider:

Axiom schema of Markov's principle:

$$MP: \quad \forall x(A(x) \lor \neg A(x)) \land \neg \neg \exists x A(x) \to \exists x A(x).$$

Axiom schema of the limited principle of omniscience:

$$LPO: \quad \forall x(A(x) \lor \neg A(x)) \to \exists xA(x) \lor \neg \exists xA(x),$$

Both principles enlarge the concept of constructivity, particularly the concept of \exists from that of intuitionistic logic H_* . However, still we have:

Theorem $H_* + MP$ and $H_* + LPO$ enjoy DP and EP. That is, MP and LPO preserve DP and EP.

Definition

A formula is said to be a Harrop-formula (H-formula) if every strictly positive subformula is neither of the form $A \lor B$ nor $\exists xA(x)$.

Theorem (Harrop)

H_{*} has the *H*(arrop)-DP and the *H*(arrop)-EP, i.e., for any *H*-formula *H*, $\mathbf{H}_* \vdash H \rightarrow A \lor B \Rightarrow \mathbf{H}_* \vdash H \rightarrow A$ or $\mathbf{H}_* \vdash H \rightarrow B$, $\mathbf{H}_* \vdash H \rightarrow \exists xA(x) \Rightarrow \mathbf{H}_* \vdash H \rightarrow A(v)$ for some *v*.

Theorem

 $\mathbf{H}_* + MP$ and $\mathbf{H}_* + LPO$ enjoy H-DP and H-EP. That is, MP and LPO preserve H-DP and H-EP.

Pointed Join of Kripke Models

Definition

- M₁, M₂: Kripke frames with the least elements 0₁ and 0₂, resp., such that the domains at 0₁ and 0₂ coincide with V(= D₁(0₁) = D₂(0₂)). A Kripke frame M is said to be the pointed join frame of M₁ and M₂, if M = {(0, V)} ↑ M₁ ⊕ M₂ with a fresh least element 0.
- (M₁, ⊨₁), (M₂, ⊨₂): Kripke models with V = D₁(0₁) = D₂(0₂). A Kripke model (M, ⊨) is said to be a pointed join model of (M₁, ⊨₁) and (M₂, ⊨₂), if M is the pointed join frame of M₁ and M₂, and the restrictions of
 - \models to \mathbf{M}_1 and \mathbf{M}_2 are \models_1 and \models_2 , resp.



Axiomatic Truth and its Preservation

Definition

A formula A is said to be axiomatically true in a Kripke model (M, \models) , if universal closures of all of substitution instances of A are true in (M, \models) .

Lemma

If A preserves its axiomatic truth in the construction of pointed join models, i.e., satisfies the following:

If A is axiomatically true in Kripke models (M₁, ⊨₁) and (M₂, ⊨₂) with V = D₁(0₁) = D₂(0₂), then A is still axiomatically true in any pointed join model of (M₁, ⊨₁) and (M₂, ⊨₂),

then $\mathbf{H}_* + A$ preserves H-DP and H-EP.

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then $\mathbf{H}_* + A$ preserves H-DP and H-EP.

Theorem

MP and LPO have this property. Hence, MP and LPO preserve H-DP and H-EP.

Another Phenomenon: Prawitz-Doorman EP

Definition

A formula is said to be a weak Harrop-formula (wH-formula) if every strictly positive subformula is not of the form $\exists xA(x)$.

Theorem (Prawitz, Doorman)

 \mathbf{H}_* has the Prawitz-Doorman EP, i.e., for any wH-formula H, $\mathbf{H}_* \vdash H \rightarrow \exists x A(x) \Rightarrow$ there exist finitely many v_1, \ldots, v_n in the vocabulary of $H \rightarrow \exists x A(x)$ such that $\mathbf{H}_* \vdash H \rightarrow A(v_1) \lor \cdots \lor A(v_n)$.

Prawitz proved EP of H_* by showing DP and the Prawitz-Doorman EP.

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Proposition

 $\mathbf{H}_* + MP$ and $\mathbf{H}_* + LPO$ fail to have the Prawitz-Doorman EP. That is, MP and LPO do not preserve the Prawitz-Doorman EP.

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In this talk, we considered preservation of DP and EP by two schemata MP and LPO in the setting of intermediate predicate logics.

	H-DP,H-EP	PD-EP
H _*	YES	YES
$\mathbf{H}_{*} + MP$, $\mathbf{H}_{*} + LPO$	YES	NO

In this talk, we considered preservation of DP and EP by two schemata MP and LPO in the setting of intermediate predicate logics.

	H-DP,H-EP	PD-EP
H _*	YES	YES
$\mathbf{H}_{*} + MP$, $\mathbf{H}_{*} + LPO$	YES	NO
$\mathbf{H}_{*} + WLPO$, $\mathbf{H}_{*} + LLPO$?	YES
$H_* + CD$	YES	NO
$H_* + WEM$	NO	YES

 $\begin{array}{l} WLPO: \ \forall x(p(x) \lor \neg p(x)) \to \neg \exists xp(x) \lor \neg \neg \exists xp(x), \\ LLPO: \ \left\{ \forall x(p(x) \lor \neg p(x)) \land \forall x(q(x) \lor \neg q(x)) \land \neg (\exists xp(x) \land \exists xq(x)) \right\} \\ \to \neg \exists xp(x) \lor \neg \exists xq(x), \\ CD: \ \forall x(p(x) \lor q) \to \forall xp(x) \lor q, \ (x \text{ is not free in } q) \\ WEM: \ \neg p \lor \neg \neg p, \end{array}$

Background Story: Ono's Problem P52 Relations betwen DP and EP in Intermediate Logics

Relations?

In intermediate predicate logics: $DP \Rightarrow EP$? $EP \Rightarrow DP$?

- (Nakamura 1983) There exists an intermediate logic having DP but lacking EP. I.e., DP ≠ EP.
- EP ⇒ DP? in intermediate logics Ono's Problem P52 (1987) (cf. Umezawa(1980), Minari(1983))

Background Story: Ono's Problem P52 Relations betwen DP and EP in Intermediate Logics

Proposition

In intermediate predicate logics, EP and DP are independent. I.e.,

- (Nakamura 1983) There exists an intermediate logic having DP but lacking EP. I.e., DP \neq EP.
- (S. 2013-15) There exists an intermediate logic having EP but lacking DP. I.e., EP ⇒ DP.

Theorem (S.2013-15)

If L is closed under the rule:

$$\frac{A \lor (p(x) \to p(y))}{A} \quad (ZR)$$

where x, y and p are distinct and do not occur in A. Then, EP of **L** implies DP of **L**.

- Do $H_* + WLPO$ and $H_* + LLPO$ have H-DP and/or H-EP?
- H-DP ⇔ DP? H-EP ⇔ EP? This problem is known as Ono's problem P54.

Remark: In intermediate propositional logic, we have: H-DP \Leftrightarrow DP.

- There must be waiting us other axiom schemata arising from constructive theories which are interesting from the viewpoint of intermediate logics!
- There must be waiting us other phenomena in intermediate logics which are interesting from the view point of constructive theories!

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