

Geometry and Effective Dimension

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Effective Dimensions

Definition

Effective Hausdorff dimension:

$$\dim_H(X) = \liminf_n \frac{C(X \upharpoonright n)}{n}.$$

Effective packing dimension:

$$\dim_\rho(X) = \limsup_n \frac{C(X \upharpoonright n)}{n}.$$

If $\dim_H X = \dim_\rho X$ we refer to X as regular and write $\dim X$.
 A and B are mutually regular if both are regular and

$$\lim_n \frac{C(A \upharpoonright n, B \upharpoonright n)}{n} \text{ exists.}$$

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A directed premetric

Definition (The d -metric)

$$d(X \rightarrow Y) := \limsup_n \frac{C(Y \upharpoonright n | X \upharpoonright n)}{n}.$$

$$d(X, Y) := \max\{d(X \rightarrow Y), d(Y \rightarrow X)\}.$$

$$X \sim_d Y := d(X, Y) = 0.$$

Theorem

d is a metric on the d -equivalence classes.

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Some Results:

- d induces a path-connected topology on the set of regular reals.
- The set of regular reals is topologically complete under d .
- The set of regular reals contains no compact neighbourhoods under d .
- If $\dim X = \dim Y$, then $d(X \rightarrow Y) = d(Y \rightarrow X)$.

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Dilutions

Definition

Let $\alpha \in [0, 1]$ and $X = x_0x_1x_2\dots$

$$\alpha X := x_0x_1\dots x_{i_1} \mathbf{000\dots 0} x_{i_1+1}x_{i_1+2}\dots x_{i_2} \mathbf{000\dots 0} x_{i_2+1}\dots$$

Where $|x_{j+1}\dots x_{j+1} \mathbf{000\dots 0}| = j + 1$, and

$$|x_{j+1}\dots x_{j+1}| = \lfloor \alpha(j + 1) \rfloor.$$

αX is the α -dilution of X .

Properties of Dilution

- 1 $0X = 0000 \dots \sim_d$ any real of dimension 0,
- 2 If X is regular, then αX is regular,
- 3 $\dim(\alpha X) = \alpha \dim(X)$,
- 4 $\alpha(\beta X) \simeq_d (\alpha\beta)X$,
- 5 $d(\alpha X \rightarrow \alpha Y) = \alpha d(X \rightarrow Y)$,
- 6 $d(\alpha X \rightarrow \beta X) = \begin{cases} 0 & \text{if } \alpha \geq \beta, \\ (\beta - \alpha) \dim X & \text{otherwise,} \end{cases}$
- 7 The map $\alpha \mapsto \alpha X$ is a continuous map from $[0, 1]$ into $2^{\mathbb{N}}$ under the d -topology.

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Compression

Theorem (Binns, Nicholson)

For every regular X of dimension α there is a unique (up to d -equivalence) Y of dimension 1, such that $X \simeq_d \alpha Y$.

Proof.

We let $X = \tau_1\tau_2\tau_3 \dots$ where $|\tau_i| = i$. Then we define

- $\gamma_1 = \tau_1$
- $\gamma_{i+1} = (\tau_{i+1} \mid \tau_1, \tau_2, \dots, \tau_i)^*$,

where the notation $(\tau_{i+1} \mid \tau_1, \tau_2, \dots, \tau_i)^*$ denotes a program of minimal length that outputs τ_{i+1} given $\tau_1, \tau_2, \tau_3, \dots, \tau_i$ as input.

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Now we let $Y = \gamma_1\gamma_2\gamma_3 \dots$



Proof cont.

Proof (very basic sketch).

Let $Y_i = \gamma_1 \gamma_2 \gamma_3 \dots \gamma_i$, and $X_i = \tau_1 \tau_2 \tau_3 \dots \tau_i$

We show that:

- $|X_i| = \mathcal{O}(i^2)$
- $|Y_i| = C(X_i) \pm \mathcal{O}(i \log i)$.
- $d(\alpha Y \rightarrow X) = 0$
- $\dim Y = 1$
- As $\dim(\alpha Y) = \alpha = \dim(X)$, this implies $d(X \rightarrow \alpha Y) = 0$ and $\alpha Y \sim_d X$.



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dim $Y = 1$ (Proof sketch)

Proof.

$$\begin{aligned} C(\gamma_{i+1} | Y_i) + C(Y_i | X_i) &\geq C(\tau_{i+1} | X_i) - o(i^2) \\ &= |\gamma_{i+1}| - o(i^2). \end{aligned}$$

But

$$\begin{aligned} C(Y_i | X_i) &= C(Y_i, X_i) - C(X_i) \pm o(i^2) \\ &= C(Y_i) - C(X_i) \pm o(i^2). \end{aligned}$$

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Randomness and Complexity

Theorem

Given any regular real X of dimension $\alpha > 0$, there is a 1-random real R such that $X \sim_d \alpha R$.

Lemma (Kučera, Gács & Merkle, Mihailović)

There is a partial computable functional Φ on $2^{\mathbb{N}}$ with the properties:

- 1 For every $Y \in 2^{\mathbb{N}}$ there is a Martin-Löf random R such that $\Phi^R = Y$.*
- 2 There is a computable function g such that $g(n)$ bounds the use of R in calculating $Y \upharpoonright n$.*
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Given X of dimension α , construct Y of dimension 1 such that $X \sim_d \alpha Y$. Now take R as in Lemma for Y . To describe $Y \upharpoonright n$ given $R \upharpoonright n$ we need only an extra $g(n) - n$ bits. Thus

$$\begin{aligned}d(R \rightarrow Y) &= \limsup_n \frac{C(Y \upharpoonright n \mid R \upharpoonright n)}{n} \\ &\leq \limsup_n \frac{g(n) - n + \mathcal{O}(1)}{n} \\ &\leq \limsup_n \frac{g(n)}{n} - 1 \\ &\leq 0\end{aligned}$$

As $\dim R = 1 = \dim Y$, we have also that $d(Y \rightarrow R) = 0$ and $Y \sim_d R$. Finally $\alpha R \sim_d \alpha Y \sim_d X$ as required. \square

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Convex combinations

We can pad a real X with bits from another real Y :

Let $r \in [0, 1]$ and $X = x_0x_1x_2\dots$, $Y = y_0y_1y_2\dots$

Define

$$r[X, Y] = x_0x_1\dots x_{i_1}y_0y_1\dots y_{j_1}x_{i_1+1}x_{i_1+2}\dots x_{i_2}y_{j_1+1}\dots y_{j_2}x_{i_2+1}\dots$$

Where $|x_{i_k+1}\dots x_{i_{k+1}}y_{j_k+1}\dots y_{j_{k+1}}| = k + 1$, and

$$i_{k+1} = \lfloor r(k + 1) \rfloor.$$

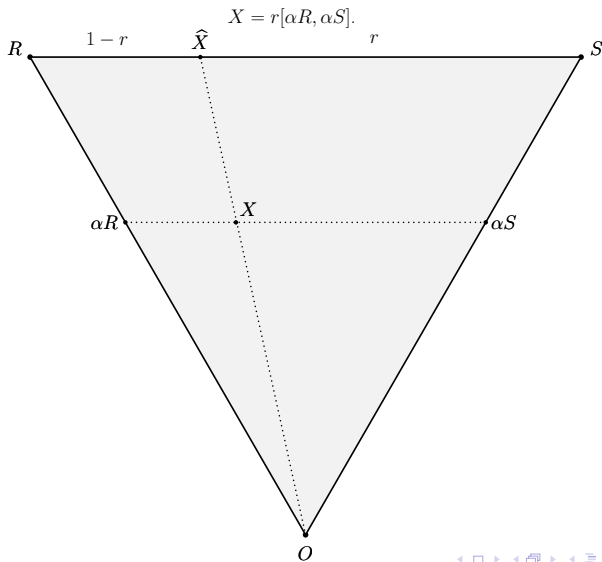
Definition

The convex hull of X and Y is

$$\mathcal{H}(X, Y) = \{r[\alpha X, \alpha Y] : r, \alpha \in [0, 1]\}$$

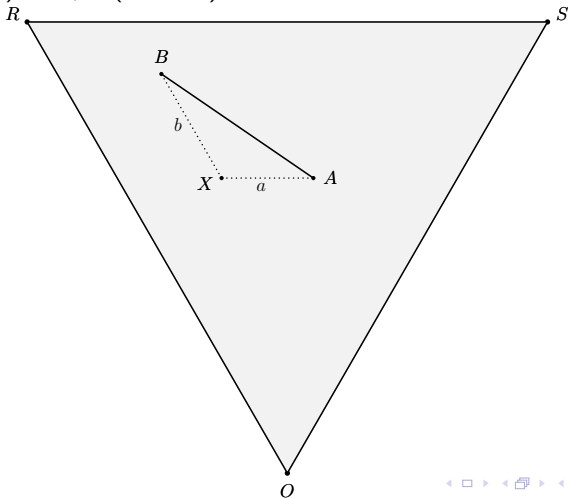
closed under d -equivalence.

The convex hull of mutually random reals R and S



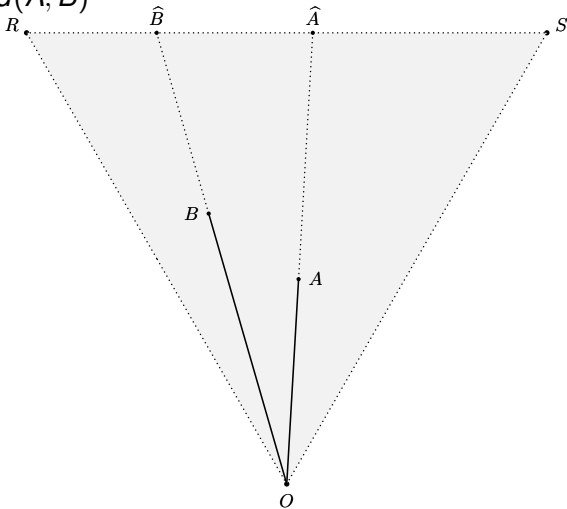
The distance function

$d(B \rightarrow X) = 0$, $d(X \rightarrow B) = b$, $d(X \rightarrow A) = d(A \rightarrow X) = a$,
 $d(B \rightarrow A) = a$, $d(A \rightarrow B) = a + b$.



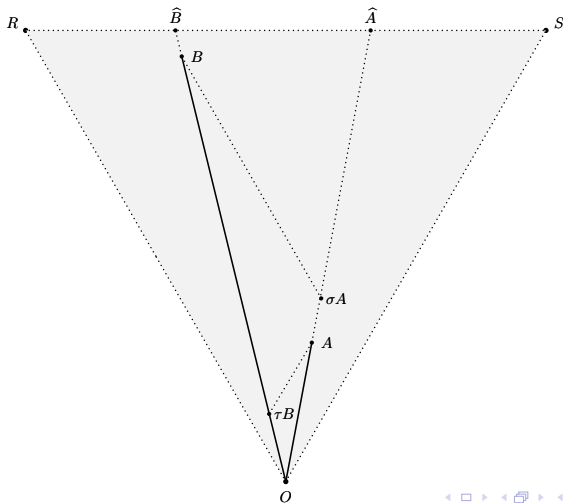
Angles

$$\angle AB = d(\widehat{A}, \widehat{B})$$



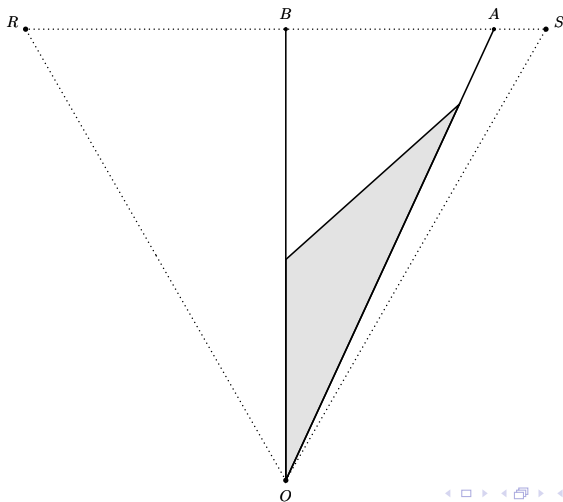
Projections: $\text{Proj}_X Y := \sup\{\alpha : d(Y \rightarrow \alpha X) = 0\}$

$\tau = \text{Proj}_B A$ $\sigma = \text{Proj}_A B$.



Other hulls: $A, B \in \mathcal{H}(R, S)$.

$$B = \frac{1}{2}[R, S], \quad A = \frac{1}{10}[R, S].$$



Question

What kind of geometry can $\mathcal{H}(A, B)$ exhibit?

The previous examples have planar hulls. They can be isometrically embedded in $\mathcal{H}(R, S)$.

Are all hulls planar? No.

Other hulls: Arbitrary mutually regular A and B

Definition

Let A and B be mutually regular reals of dimension 1. Then A and B form a coherent pair if

$$\angle AB = \frac{(1 - \sigma)(1 - \tau)}{1 - \sigma\tau},$$

where $\sigma = \text{Proj}_A B$ and $\tau = \text{Proj}_B A$.

Theorem

$\mathcal{H}(A, B)$ is planar if and only if A and B are coherent.

Theorem

Not all mutually regular reals are coherent.

Other hulls: Arbitrary mutually regular A and B

Definition

Let A and B be mutually regular reals of dimensions a and b respectively. Then A and B form a coherent pair if

$$\angle AB = \frac{(b - a\sigma)(a - b\tau)}{ab(1 - \sigma\tau)},$$

where $\sigma = \text{Proj}_A B$ and $\tau = \text{Proj}_B A$.

Theorem

$\mathcal{H}(A, B)$ is planar if and only if A and B are coherent.

Theorem

Not all mutually regular reals are coherent.

Not all $A B$ are coherent

Proof.

Let $R = r_1 r_2 r_3 r_4 \dots r_n \dots$ be a random real. Let

$$A = r_0 r_2 r_4 \dots r_{2n} \dots$$

$$B = r_0 r_3 r_6 \dots r_{3n} \dots$$

Both A and B are random and so are dimension 1. But

- $\angle AB = d(B \rightarrow A) = 2/3$
- $\text{Proj}_B A = 0$
- $\text{Proj}_A B = 0$
- $2/3 \neq \frac{(1-0)(1-0)}{1-0} = 1.$



Other directions

- Given a coherent pair A and B , does there exist a pair of mutually random reals R, S such that

$$\mathcal{H}(A, B) \subseteq \mathcal{H}(R, S)?$$

- Given a mutually regular pair A, B is there a curve of length $d(A, B)$ connecting A and B ?
- Given a mutually regular pair A and B , does there exist a (unique?) C of minimal dimension such that $d(B, C \rightarrow A) = 0$? Where

$$d(B, C \rightarrow A) := \limsup_n \frac{C(A \upharpoonright n \mid B \upharpoonright n \oplus C \upharpoonright n)}{n}.$$

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